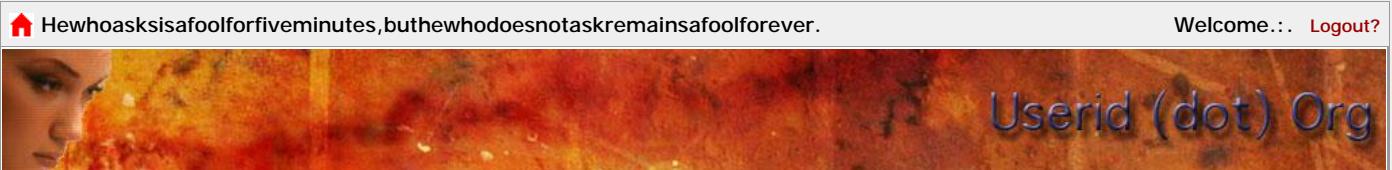


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Chapter3. Enhanced **IPSec** Features

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GRE and IPSec

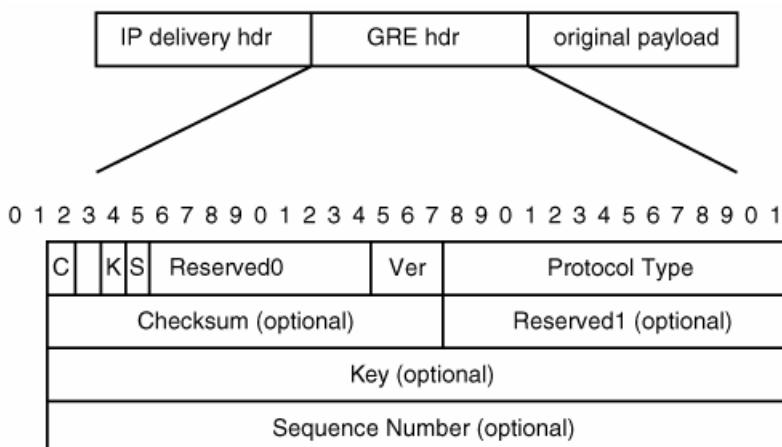
Designing a VPN using **IPSec** for connectivity between peers has inherent limitations. These are:

- **IPSec** can encrypt/decrypt only IP traffic.
- IP traffic destined to a multicast or broadcast IP address cannot be handled by **IPSec**, which means that IP multicast traffic cannot traverse the **IPSec tunnel**. Also, many routing protocols (such as EIGRP, OSPF, and RIPv2) use a multicast or broadcast address; therefore, dynamic routing using these routing protocols cannot be configured between **IPSec** peers.

[/]
 [tr]
 [wired]
 [bbc]
 [fark]
 [salon]
 [gn]
 [sb]
 [jmro]
 [g&m]
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 [ud]
 [cbc]
 [istop]
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These limitations can be overcome by configuring an IP-encapsulated **GRE tunnel** between the peers and applying **IPSec** protection on the **GRE/IP tunnel**. RFC 2784 covers **GRE** in detail. It essentially **GRE**-encapsulates any payload in an IP unicast packet destined to the **GRE** endpoint. A **GRE**-encapsulated packet is shown in [Figure 3-7](#).

Figure 3-7. **GRE-Encapsulated Packet**



When **GRE** is used in conjunction with **IPSec**, either **tunnel** mode or transport mode can be used. **Tunnel** mode adds an **IPSec** IP header to the **GRE** packet whereas **IPSec** transport mode uses the original **GRE** packet's IP header. Figures 3-8 and 3-9 show **GRE** with **IPSec** in **tunnel** mode and transport mode, respectively.

Figure 3-8. GRE-Encapsulated Packet in IPSec Tunnel Mode

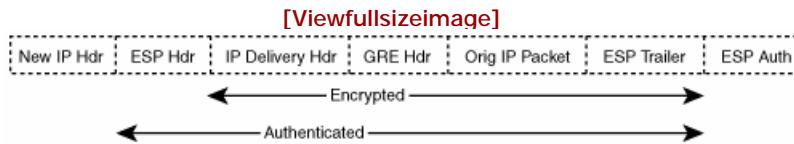
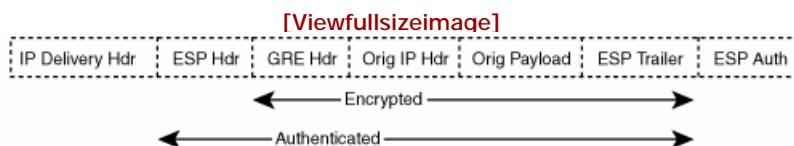


Figure 3-9. GRE-Encapsulated Packet in IPSec Transport Mode



IPSec transport mode is the most efficient way to combine **GRE** and **IPSec** together because **GRE** encapsulation already places a new IP header on the payload. The use of **IPSec** transport mode, however, requires that the **GRE** encapsulation use an IP source and destination address that is reachable via the IP path to its peer. **GRE** adds 24 bytes of overhead to the original IP packet. In conjunction with **IPSec** transport mode, **GRE** encapsulation adds 4 bytes of extra overhead in comparison to the 20 bytes of overhead added by **IPSec** tunnel mode. Although the additional overhead and extra processing for **GRE** encapsulation reduce the overall throughput and may impact the latency of the connection, the benefits of using **GRE** with **IPSec** far outweigh the impact.

Use of **GRE** with **IPSec** also has the useful side effect of making **IPSec** VPN configurations simpler. Traditional **IPSec** configuration between peers requires **IPSec** policy configuration to specify the protected subnets so that traffic destined to the protected subnets is encrypted or decrypted. Everytime a new protected subnet is added or deleted, the **IPSec** policy configuration needs to be updated on both peers. Also, the security policy database (SPD) and security association database (SADB) size can get quite large depending on the **IPSec** SA bundles negotiated and installed. This usually has an impact on the overall performance and scalability of the security gateway. Using **GRE** with **IPSec** significantly reduces the configuration complexity, as the policy in the SPD need to match the traffic only to the **GRE** endpoint addresses. The addition or deletion of the protected subnets requires no change in the configuration of the **IPSec** SPD peers. The protected subnet traffic is directed to be encrypted by being routed out the **GRE tunnel** interface (everything that goes through the **GRE tunnel** will be encrypted). This does mean that the granularity for what gets encrypted is now at the IP address level rather than the port level in the transport header. In general, this is not an issue because usually all traffic for hosts on protected subnets is to be encrypted.

The **GRE** encapsulation does increase the size of the packet and potentially causes fragmentation issues. Packet fragmentation can be avoided by ensuring that PMTUD is enabled. To ensure that PMTUD works as expected, ICMP code 3 type 4 messages must be allowed through the network. If ICMP code 3 type 4 messages are not supported in the network, setting a lower MTU size on the **tunnel** interface will cause fragmentation before encryption to happen, achieving the same effect as Look Ahead Fragmentation. If the end hosts support PMTUD, then they will match the packet size to the configured MTU. Both the **IPSec** and **GRE** specification support Path MTU Discovery by allowing the copy of the DF bit of the original IP header into the newly built IP header. This is usually a configuration option of the devices.

If the end host does not support PMTU and the DF bit is not set, the packet will be fragmented and sent over the **GRE-encrypted tunnel**. One benefit to reducing the MTU on the **GRE tunnel** is that packet fragmentation occurs before it hits the encryption process; therefore, there assembly is done on the end host and not on the peer **IPSec** gateway.

Example3-10 shows the configuration of **GRE** tunnels protected by **IPSec** on SPOKE-1-EAST. Note the **tunnel protection ipsec profile gre** command under the **GRE tunnel** interface configuration, which is a new way of protecting **GRE** tunnels using **IPSec**, wherein the physical interface that the **GRE tunnel** traverses does not need the crypto map configuration. Compare this configuration with **Example3-11**, which shows the old way of **IPSec**-protected **GRE** in IOS.

Example3-10. **GRE** and **IPSec** Using **Tunnel** Protection CLI

```
spoke-1-east

!
crypto isakmp policy 1
  authentication pre-share
crypto isakmp key cisco address 9.1.1.35
crypto isakmp keepalive 60 2

!
crypto ipsec transform-set test esp-3des esp-sha-hmac
!
crypto ipsec profile gre
  set transform-set test
!
int tu0
  ip address 192.168.1.1 255.255.255.0
  ip mtu 1400
  ip tcp adjust-mss 1360
  tunnel source s0/0
  tunnel destination 9.1.1.35
  tunnel path-mtu-discovery
  tunnel protection ipsec profile gre

!
interface Serial0/0
  ip address 9.1.1.146 255.255.255.252
  crypto map vpn
!
interface Ethernet0/1
  ip address 10.0.68.1 255.255.255.0
  half-duplex
!
ip classless
ip route 0.0.0.0 0.0.0.0 9.1.1.145
ip route 10.0.1.0 0.0.0.255 tu0
```

Example3-11. **GRE Tunnel** Keepalive

```
spoke-1-east

!
crypto isakmp policy 1
  authentication pre-share
crypto isakmp key cisco address 9.1.1.35
crypto isakmp keepalive 60 2
```

```

!
crypto ipsec transform-set test esp-3des esp-sha-hmac
mode transport
!
crypto map vpn 1 ipsec-isakmp
  set peer 9.1.1.35
  set transform-set test
  match address 100

!
int tunnel0
ip address 192.168.1.1 255.255.255.0
ip mtu 1400
ip tcp adjust-mss 1360
tunnel path-mtu-discovery
tunnel source s0/0
tunnel destination 9.1.1.35
keepalive 20 3

!
interface Serial0/0
  ip address 9.1.1.146 255.255.255.252
  crypto map vpn

!
interface Ethernet0/1
  ip address 10.0.68.1 255.255.255.0
  half-duplex
!
ip classless
ip route 0.0.0.0 0.0.0.0 9.1.1.145
ip route 10.0.1.0 0.0.0.255 tu0
!
access-list 100 permit gre host 9.1.1.146 host 9.1.1.35

spoke-1-east#show cry isa sa
dst          src          state      conn-id    slot
9.1.1.35    9.1.1.146   QM_IDLE   82         0

spoke-1-east#show int tu0
Tunnel0 is up, line protocol is up
  Hardware is Tunnel
  Interface is unnumbered. Using address of Ethernet0/1 (10.0.68.1)
  MTU 1514 bytes, BW 9 Kbit, DLY 500000 usec,
    reliability 255/255, txload 1/255, rxload 1/255
  Encapsulation TUNNEL, loopback not set
  Keepalive set (20 sec), retries 3

Tunnel source 9.1.1.146 (Serial0/0), destination 9.1.1.35
Tunnel protocol/transport GRE/IP, key disabled, sequencing disabled
Tunnel TTL 255
Checksumming of packets disabled, fast tunneling enabled
Last input 00:00:07, output 00:00:07, output hang never
Last clearing of "show interface" counters never
Input queue: 0/75/0/0 (size/max/drops/flushes); Total output drops:14
Queueing strategy: fifo
Output queue: 0/0 (size/max)
5 minute input rate 0 bits/sec, 0 packets/sec
5 minute output rate 0 bits/sec, 0 packets/sec
  483 packets input, 36396 bytes, 0 no buffer
  Received 0 broadcasts, 0 runts, 0 giants, 0 throttles

```

```

0 input errors, 0 CRC, 0 frame, 0 overrun, 0 ignored, 0 abort
168 packets output, 8596 bytes, 0 underruns
0 output errors, 0 collisions, 0 interface resets
0 output buffer failures, 0 output buffers swapped out

Tunnel0: sending keepalive, 9.1.1.35->9.1.1.146 (len=24 ttl=255),
counter=1
Tunnel0: GRE/IP encapsulated 9.1.1.146->9.1.1.35 (linktype=7, len=48)
IP: s=9.1.1.146 (Tunnel0), d=9.1.1.35 (Serial0/0), len 48, sending,
proto =47
IP: s=9.1.1.35 (Serial0/0), d=9.1.1.146 (Serial0/0), len 24, rcvd 3,
proto=47
Tunnel0: GRE/IP to decaps 9.1.1.35->9.1.1.146 (len=24 ttl=253)
Tunnel0: keepalive received, 9.1.1.35->9.1.1.146 (len=24 ttl=253),
reset

```

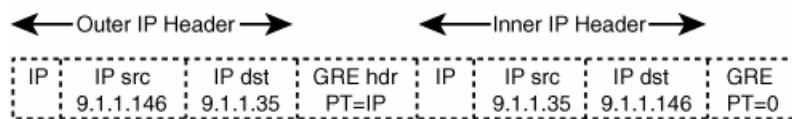
The primary motivation for using **GRE** with **IPSec** is its ability to run dynamic routing protocols such as OSPF or EIGRP between sites for advertising the protected subnets. Dynamic routing also implicitly helps in failover situations. On the other hand, if static routes are used for the reachability of the protected subnets with **GRE**, then there is no way for the **IPSec** peers to know that the protected subnets are not reachable when the **GRE tunnel** endpoints are not reachable anymore. To facilitate knowledge of the **tunnel** status, **GRE** keepalive messages can be configured on the **GRE** peers. [Example 3-11](#) shows configuration of **GRE** keepalive under the **tunnel** interface. **GRE tunnel** comes up as soon as it sees the default route in the routing table on SPOKE-1-EAST. Now, if VPN-GW1-EAST goes down, you want the **tunnel** interface on SPOKE-1-EAST to go down—this is facilitated by **GRE** keepalives.

Note

GRE keepalives are not supported with the **tunnel protection IPSec** configurations syntax. Therefore, for **GRE** keepalive functionality, you need to use the old-style cryptomap configuration.

[Figure 3-10](#) shows the format of a **GRE** keepalive packet from SPOKE-1-EAST's perspective.

[Figure 3-10. GRE Keepalive Packet](#)



Notice that the destination IP address in the inner IP header is SPOKE-1-EAST's **tunnel** source address (9.1.1.146) and the source IP address in the inner IP header is that of the **tunnel** destination address of VPN-GW1-EAST (9.1.1.35). The protocol type (PT) in the inner header is set to 0. This payload is sent in a **GRE tunnel** with a protocol type of IP in the outer header with source and destination addresses as configured on the **tunnel** interfaces shown in [Example 3-11](#). The **tunnel** keepalive counter is incremented by one as shown in the debug snapshots in [Example 3-11](#). As the **GRE tunnel** is protected via the cryptomap, this keepalive packet will be encrypted when it leaves SPOKE-1-EAST. The packet will reach the far end **tunnel** endpoint peer (VPN-GW1-EAST) via normal routing. Upon arrival on VPN-GW1-EAST, the packet will get decrypted and then decapsulated; the resulting packet will have a source IP address of VPN-GW1-EAST and a destination IP address of SPOKE-1-EAST. This packet will now make its way back to SPOKE-1-EAST through the **GRE tunnel**.

which is again encrypted. The packet, on reaching SPOKE-1-EAST after decryption and GRE decapsulation, will result in a PT of 0, which will signify that this is a keepalive packet that it originally constructed. The receipt of this packet signifies the GRE tunnel is up. The tunnel keepalive counter will be reset to 0 and the packet will be discarded. This process is repeated periodically by each GRE peer.

When VPN-GW1-EAST becomes unreachable from SPOKE-1-EAST for whatever reason, SPOKE-1-EAST will continue to construct and send the keepalive packets as well as normal traffic. Because the keepalives do not come back, the tunnel line protocol will stay up as long as the tunnel keepalive counter is less than the configured retry value. When the number of retries exceeds the configured value, the line protocol will be brought down on the tunnel interface on VPN-SPOKE1-EAST.

In the up/down state, the tunnel will not forward or process any traffic apart from the keepalive packets. The reception of a keepalive packet from VPN-GW1-EAST would imply that the tunnel endpoint is again reachable; when this happens, the tunnel keepalive counter will be reset to 0 and the line protocol will change state back to up, resuming data traffic.

One thing worth mentioning is that GRE keepalive packets are sent out with the TOS bit set to 5 so that the routers process those packets with higher priority, even during congestion.

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